

Knowledge Representation for the Semantic Web

Winter Quarter 2010

Slides 6 – 02/04/2010

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Slides are based on



Pascal Hitzler, Markus Krötzsch, Sebastian Rudolph

Foundations of Semantic Web Technologies

Chapman & Hall/CRC, 2010

Flyer with special offer is available.



Pascal Hitzler Markus Krötzsch Sebastian Rudolph

CRC Press operations

http://www.semantic-web-book.org



Today: Description Logics







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1. Basic Ideas

- 2. The Description Logic SROIQ(D)
- 3. Different Perspectives
- 4. Class Project
- 5. Class Presentations





- individuals (written as URIs)
 - also: constants (FOL), resources (RDF), instances
 - http://example.org/sebastianRudolph
 - http://www.semantic-web-book.org/
 - we write these lowercase and abbreviated, e.g.
 "sebastianRudolph"
- classes (also written as URIs!)
 - also: concepts, unary predicates (FOL)
 - we write these uppercase, e.g. "Father"
- properties (also written as URIs!)
 - also: roles (DL), binary predicates (FOL)
 - we write these lowercase, e.g. "hasDaughter"



DL syntax

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Woman \sqsubseteq Person





- Person(mary) :mary rdf:type :Person .
 - :Woman rdfs:subClassOf :Person .
 - Person ≡ HumanBeing (class equivalence):
 Person ⊑ HumanBeing AND HumanBeing ⊑ Person
- hasWife(john,mary)
 :john :hasWife :mary .
- hasWife
 hasSpouse

 hasWife rdfs:subPropertyOf :hasSpouse
 - hasSpouse \equiv marriedWith (class equivalence)



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DL syntax

FOL syntax



• Person(mary) • Person(mary)

ABox statements

- Woman \sqsubseteq Person $\forall x (Woman(x) \rightarrow Person(x))$
 - Person \equiv HumanBeing (class equivalence)
- hasWife(john,mary)
 hasWife(john,mary)
 - hasWife \sqsubseteq hasSpouse $\forall x \forall y \text{ (hasWife}(x,y) \rightarrow hasSpouse(x,y))$
 - hasSpouse \equiv marriedWith (class equivalence)

TBox statements



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Special classes and properties



- owl:Thing (RDF syntax)
 - DL-syntax: ⊤
 - contains everything
- owl:Nothing (RDF syntax)
 - DL-syntax: \perp
 - empty class
- owl:topProperty (RDF syntax)
 - DL-syntax: U
 - every pair is in U
- owl:bottomProperty (RDF syntax)
 - empty property



Class constructors

• conjunction

- Mother \equiv Woman \sqcap Parent

"Mothers are exactly those who are women and parents."

• disjunction

 $\forall x (Parent(x) \leftrightarrow Mother(x) \lor Father(x))$

 $\forall x (Mother(x) \leftrightarrow Woman(x) \land Parent(x))$

- Parent \equiv Mother \sqcup Father

"Parents are exactly those who are mothers or fathers."

• negation

 $\forall x \text{ (ChildlessPerson(x)} \leftrightarrow \text{Person(x)} \land \neg \text{Parent(x))}$

- ChildlessPerson \equiv Person $\sqcap \neg$ Parent

"ChildlessPersons are exactly those who are persons and who are not parents."





Class constructors

- existential quantification
 - only to be used with a role also called a *property* restriction $\forall x (Parent(x) \leftrightarrow$
 - Parent $\equiv \exists$ hasChild.Person "Parents are exactly those who have at least one child which is a Person."
- universal quantification
 - only to be used with a role also called a property restriction
 - Person □ Happy ≡ ∀hasChild.Happy "A (person which is also happy) is exactly (something all children of which are happy)."

 $\begin{array}{l} \forall x \ (\text{Person}(x) \land \text{Happy}(x) \leftrightarrow \\ \forall y \ (\text{hasChild}(x,y) \rightarrow \text{Happy}(y))) \end{array} \end{array}$

Class constructors can be nested arbitrarily

 $\exists y (hasChild(x,y) \land Person(y)))$





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The description logic ALC

- ABox expressions: Individual assignments Property assignments
- TBox expressions subclass relationships
 - conjunction disjunction negation

property restrictions

Complexity: ExpTime

Father(john) hasWife(john,mary)

 \equiv for equivalence

Also: \top , \bot



 \square

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The Description Logic ALC



- Set of *individuals* a,b,c,...
- Set of atomic classes (class names) A,B,...
- Set of role names R,S,...
- (Complex) class expressions are constructed as:

 $C,D::=A\mid\top\mid\perp\mid\neg C\mid C\sqcap D\mid C\sqcup D\mid\forall R.C\mid\exists R.C$

- A *TBox* is a set of statements of the form $C \equiv D$ or $C \sqsubseteq D$, where C and D are class expressions. They are called *general inclusion axioms*.
- An ABox consists of statements of the form C(a) or R(a,b), where C is a class expression, R is a role, and a, b are individuals.





Human ⊑ ∃hasParent.Human Orphan ⊑ Human □ ∀hasParent.¬Alive Orphan(harrypotter) hasParent(harrypotter,jamespotter)





- ALC + role chains = SR
- hasParent \circ hasBrother \sqsubseteq hasUncle.

 $\forall x \; \forall y \; (\exists z \; ((hasParent(x,z) \land hasBrother(z,y)) \rightarrow hasUncle(x,y)))$

includes top property and bottom property

- includes S = ALC + transitivity
 - hasAncestor o hasAncestor
 LasAncestor
- includes SH = S + role hierarchies
 - hasFather ⊑ hasParent



Understanding SROIQ(D)



- O nominals (closed classes)
 - MyBirthdayGuests \equiv {bill,john,mary}
 - Note the difference to MyBirthdayGuests(bill) MyBirthdayGuests(john) MyBirthdayGuests(mary)
- Individual equality and inequality (no unique name assumption!)
 - bill = john
 - ${bill} \equiv {john}$
 - bill ≠ john
 - {bill} \sqcap {john} $\equiv \bot$



Understanding SROIQ(D)

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- I inverse roles
 - hasParent \equiv hasChild
 - Orphan $\equiv \forall hasChild$.Dead
- Q qualified cardinality restrictions
 - <4 hasChild.Parent(john)</p>
 - HappyFather $\equiv \geq$ 2 hasChild.Female
 - Car ⊑ =4hasTyre.⊤
- Complexity SHIQ, SHOQ, SHIO: ExpTime. Complexity SHOIQ: NExpTime Complexity SROIQ: N2ExpTime





Properties can be declared to be

- Transitive
- Symmetric hasSpouse
- Asymmetric
- Reflexive
- Irreflexive parentOf
- Functional hasHusband
- InverseFunctional hasHusband

 $\begin{array}{l} \mathsf{R}(a,b) \text{ and } \mathsf{R}(b,c) \to \mathsf{R}(a,c) \\ \mathsf{R}(a,b) \to \mathsf{R}(b,a) \\ \mathsf{R}(a,b) \to \mathsf{not} \ \mathsf{R}(b,a) \\ \mathsf{R}(a,a) \text{ for all } a \\ \mathsf{not} \ \mathsf{R}(a,a) \text{ for any } a \\ \mathsf{R}(a,b) \text{ and } \mathsf{R}(a,c) \to \mathsf{b=c} \\ \mathsf{R}(a,b) \text{ and } \mathsf{R}(c,b) \to \mathsf{a=c} \end{array}$

called property characteristics



hasAncestor

hasChild

hasRelative



(D) – datatypes

- so far, we have only seen properties with individuals in second argument, called object properties or abstract roles (DL)
- properties with datatype literals in second argument are called data properties or concrete roles (DL)
- In OWL allowed are many XML Schema datatypes, including xsd:integer, xsd:string, xsd:float, xsd:booelan, xsd:anyURI, xsd:dateTime

and also e.g. owl:real





(D) – datatypes

- hasAge(john, "51"^^xsd:integer)
- additional use of *constraining facets* (from XML Schema)
 - e.g. Teenager = Person $\sqcap \exists$ hasAge.(xsd:integer: ≥12 and ≤19)

note: this is not standard DL notation! It's really only used in OWL.





further expressive features

- Self
 - PersonCommittingSuicide $\equiv \exists$ kills.Self
- Keys (not really in SROIQ(D), but in OWL)
 - set of (object or data) properties whose values uniquely identify an object
- disjoint properties
 - Disjoint(hasParent,hasChild)
- explicit anonymous individuals
 - as in RDF: can be used instead of named individuals





- ABox assignments of individuals to classes or properties
- ALC: <u>□</u>, ≡ for classes
 □, □, ¬, ∃, ∀
 ⊤, ⊥
- SR: + property chains, property characteristics, role hierarchies ⊑
- SRO: + nominals {o}
- SROI: + inverse properties
- SROIQ: + qualified cardinality constraints
- SROIQ(D): + datatypes (including facets)
- + top and bottom roles (for objects and datatypes)
- + disjoint properties
- + Self
- + Keys (not in SROIQ(D), but in OWL)





Available in OWL (see later) as syntactic sugar for DL axioms.

- disjoint classes
 - Apple \sqcap Pear $\sqsubseteq \bot$
- disjoint union
 - Parent = Mother ⊔ Father Mother \sqcap Father $\sqsubseteq \bot$
- negative property assignments (also for datatypes)
 ¬hasAge(jack,"53"^^xsd:integer)



Е.sis

arbitrary property chain axioms lead to undecidability

- **restriction**: set of property chain axioms has to be *regular*
 - there must be a strict linear order < on the properties</p>
 - every property chain axiom has to have one of the following forms:
 - $R \circ R \sqsubseteq R$ $S^- \sqsubseteq R$ $S_1 \circ S_2 \circ _ \circ S_n \sqsubseteq R$ $R \circ S_1 \circ S_2 \circ _ \circ S_n \sqsubseteq R$ $S_1 \circ S_2 \circ _ \circ S_n \circ S_n \subseteq R$ $S_1 \circ S_2 \circ _ \circ S_n \circ S_n \subseteq R$ $S_1 \circ S_2 \circ _ \circ S_n \circ S_n \circ R \subseteq R$

Example 1: $R \circ S \sqsubseteq R$ $S \circ S \sqsubseteq S$ $R \circ S \circ R \sqsubseteq T$

 \rightarrow regular with order S < R < T

Example 2: $\mathbf{R} \circ \mathbf{T} \circ \mathbf{S} \sqsubseteq \mathbf{T}$

 \rightarrow not regular because form not admissible

Example 3: $R \circ S \sqsubseteq S \circ R \sqsubseteq R$

 \rightarrow not regular because no adequate order exists





- combining property chain axioms and cardinality constraints may lead to undecidability
- restriction: use only *simple* properties in cardinality expressions (i.e. those which cannot be – directly or indirectly – inferred from property chains)
- technically:
 - for any property chain axiom $S_1 \circ S_2 \circ _$ o $S_n \sqsubseteq R$ with n>1, R is non-simple
 - for any subproperty axiom S ⊑ R with S non-simple, R is non-simple
 - all other properties are simple
- **Example:** $Q \circ P \sqsubseteq R$ $R \circ P \sqsubseteq R$ $R \sqsubseteq S$ $P \sqsubseteq R$ $Q \sqsubseteq S$ non-simple: R, S simple: P, Q



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OWL – Extralogical Features

- OWL ontologies have URIs and can be referenced by others via

 import statements
- Namespace declarations
- Entity declarations (must be done)
- Versioning information etc.
- Annotations
 - Entities and axioms (statements) can be endowed with annotations, e.g. using rdfs:comment.
 - OWL syntax provides annotation properties for this purpose.

Note: We still have to give a syntax for OWL – forthcoming.





The modal logic perspective



- Description logics can be understood from a modal logic perspective.
- Each pair of ∀R and ∃R statements give rise to a pair of modalities.
- Essentially, some description logics are multi-modal logics.



The RDFS perspective



RDFS semantics is weaker

- :mary rdf:type :Person .
- :Mother rdfs:subClassOf :Woman .
- :john :hasWife :Mary .
- :hasWife rdfs:subPropertyOf :hasSpouse
- :hasWife rdfs:range :Woman .
- :hasWife rdfs:domain :Man .

- Person(mary)
- Mother 🗆 Woman
- hasWife(john,mary)
- hasWife ⊑ hasSpouse

- ⊤ ⊑ ∀hasWife.Woman
- ⊤ ⊑ ∀hasWife .Man or ∃hasWife.⊤ ⊑ Man

RDFS also allows to

- make statements about statements \rightarrow only possible through annotations in OWL (not present in SROID(D))
- mix class names, individual names, property names (they are all URIs) $\rightarrow punning \text{ in OWL}$



Punning



- Description logics impose *type separation*, i.e. names of individuals, classes, and properties must be disjoint.
- In OWL 2 Full, type separation does not apply.
- In OWL 2 DL, type separation is relaxed, but a class X and an individual X are interpreted semantically as if they were different.
- Father(john) SocialRole(Father)
- See further below on the two different types/semantics for OWL: OWL DL and OWL Full.



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• none this time.



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Class presentations – scheduled



- RDFa embedding RDF in HTML (W3C standard) Pavan, Thursday 28th of January
- Scalable Distributed Reasoning using MapReduce (Urbani, Kotoulas, Oren, van Harmelen, ISWC2009)
 Wenbo, Thursday 28th of January

All remaining presentations will be in the last week

- Semantic MediaWiki, Vinh, to be scheduled
- Linked Open Data, Ashutosh, to be scheduled
- FOAF, Hemant, to be scheduled
- Virtuoso, Pramod, to be scheduled
- Prateek, Conjunctive Queries for OWL
- Raghava, DL-Lite





Thursday 4th of February: OWL Part 1 Tuesday 9th of February: OWL Part 2 Thursday 11th of February: OWL Part 3 Tuesday 23rd of February: Exercise Session Thursday 25th of February: OWL Part 4 Week from March 8th: Class Presentations Friday March 12th: most exams

Estimated breakdown of sessions:

Intro + XML: 2	RDF: 3.3
OWL: 4	SPARQL: 1
Class Project Session: 2	Class Presentations: 3
Exercise sessions: 2.7	

