

Knowledge Representation for the Semantic Web

Winter Quarter 2010

Slides 11 - 03/08/2010

Pascal Hitzler

010 Kno.e.sis Center Wright State University, Dayton, OH

http://www.knoesis.org/pascal/





Slides are based on



Pascal Hitzler, Markus Krötzsch, Sebastian Rudolph

Foundations of Semantic Web Technologies

Chapman & Hall/CRC, 2010

Flyer with special offer is available.



Pascal Hitzler Markus Krötzsch Sebastian Rudolph

CRC Press operations

http://www.semantic-web-book.org



Today: Wrap-Up



• Main Messages

- What Is Semantics Revisited
- OWL At Its Expressive Limits
- A Very Personal Semantic Web History
- Making OWL Fit For Practice





- How to model in RDF and OWL
- What is model-theoretic semantics
- How to compute logical consequences in RDF and OWL



Today: Wrap-Up



- Main Messages
- What Is Semantics Revisited
- OWL At Its Expressive Limits
- A Very Personal Semantic Web History
- Making OWL Fit For Practice



What Is Semantic Web Semantics?



- Opinions Differ. Here's my take.
- Semantic Web requires a *computable* semantics.
- I.e., the semantics must be a formal entity which is clearly defined and automatically computable.
- Ontology languages provide this by means of their formal semantics.
- Semantic Web Semantics is given by a relation the *logical* consequence relation.
- Note: This is considerably more than saying that the semantics of an ontology is the set of its logical consequences!



Today: Wrap-Up



- Main Messages
- What Is Semantics Revisited
- OWL At Its Expressive Limits
- A Very Personal Semantic Web History
- Making OWL Fit For Practice



OWL At Its Expressive Limits



- There are a lot of things that cannot be said in OWL.
- We will talk about a few such things and general ideas how to address them.



Simple Reasoning









Less Simple Reasoning







Sophisticated Application Needs









Rules are often considered an intuitive form of knowledge representation

- Man(x) \land hasBrother(x,y) \land hasChild(y,z) \rightarrow Uncle(x)
 - Man \sqcap ∃hasBrother.∃hasChild. \top \sqsubseteq Uncle
- ThaiCurry(x) → ∃contains.FishProduct(x)
 - ThaiCurry ⊑ ∃contains.FishProduct
- kills(x,x) \rightarrow suicide(x) suicide(x) \rightarrow kills(x,x) - \exists kills.Self \sqsubseteq suicide suicide \sqsubseteq \exists kills.Self

Note: with these two axioms,

suicide is basically the same as kills



Some things you can say in OWL



- NutAllergic(x) \land NutProduct(y) \rightarrow dislikes(x,y)
 - NutAllergic ≡ ∃nutAllergic.Self
 NutProduct ≡ ∃nutProduct.Self
 nutAllergic ∘ U ∘ nutProduct ⊑ dislikes
- dislikes(x,z) ∧ Dish(y) ∧ contains(y,z) → dislikes(x,y)
 - Dish ≡ ∃dish.Self
 dislikes ∘ contains⁻ ∘ dish ⊑ dislikes
- worksAt(x,y) ∧ University(y) ∧ supervises(x,z) ∧ PhDStudent(z) → professorOf(x,z)
 - ∃worksAt.University ≡ ∃worksAtUniversity.Self
 PhDStudent ≡ ∃phDStudent.Self
 worksAtUniversity ∘ supervises ∘ phDStudent ⊑ professorOf



DL Rules: definition

- Tree-shaped bodies
- First argument of the conclusion is the root
- $C(x) \land R(x,a) \land S(x,y) \land D(y) \land T(y,a) \rightarrow E(x)$ - $C \sqcap \exists R.\{a\} \sqcap \exists S.(D \sqcap \exists T.\{a\}) \sqsubseteq E$



duplicating nominals is ok







DL Rules: definition

- Tree-shaped bodies
- First argument of the conclusion is the root
- $C(x) \land R(x,a) \land S(x,y) \land D(y) \land T(y,a) \rightarrow V(x,y)$

 $C \sqcap \exists R.\{a\} \sqsubseteq \exists R1.Self$ $D \sqcap \exists T.\{a\} \sqsubseteq \exists R2.Self$ $R1 \circ S \circ R2 \sqsubseteq V$







DL Rules: definition



- Tree-shaped bodies
- First argument of the conclusion is the root
- complex classes are allowed in the rules
 - Mouse(x) $\land \exists$ hasNose.TrunkLike(y) \rightarrow smallerThan(x,y)
 - ThaiCurry(x) $\rightarrow \exists$ contains.FishProduct(x)

Note: This allows to reason with unknowns (unlike rules)

 allowed class constructors depend on the chosen underlying description logic!

SROIQ Rules can be transformed back into **SROIQ**!





 $extsf{Pair}(p) \land extsf{consistOf}(p, x) \land extsf{consistOf}(p, y) \land extsf{differentFrom}(x, y) \land extsf{River}(r) \land extsf{inBetween}(r, p) \land extsf{rightBankOf}(x, r) \rightarrow extsf{leftBankOf}(y, r).$

- Cannot be expressed in SROIQ (is not a SROIQ Rule).
- Extending OWL with such more general rules leads to undecidability.

[Example due to Dong-Po Deng, presented at GeoS2009]







 $extsf{Pair}(p) \land extsf{consistOf}(p, x) \land extsf{consistOf}(p, y) \land extsf{differentFrom}(x, y) \land extsf{River}(r) \land extsf{inBetween}(r, p) \land extsf{rightBankOf}(x, r) \rightarrow extsf{leftBankOf}(y, r).$

• Read rule as a first-order predicate logic formula.

Semantically okay, but leads to undedicability in combination with OWL.





 $extsf{Pair}(p) \land extsf{consistOf}(p, x) \land extsf{consistOf}(p, y) \land extsf{differentFrom}(x, y) \land extsf{River}(r) \land extsf{inBetween}(r, p) \land extsf{rightBankOf}(x, r) \rightarrow extsf{leftBankOf}(y, r).$

- Semantically restrict rule, such that it applies only to individuals which are explicitly contained in the knowledge base.
 I.e., those with known URIs.
- DL-safe SWRL combined with OWL is decidable.
- Formalism supported, e.g., by Pellet.





ThaiCurry ⊑ ∃contains.{peanutOil} ⊤ ⊑ ∀orderedDish.Dish

$$\begin{split} & \mathsf{NutAllergic}(x) \land \mathsf{NutProduct}(y) \to \mathsf{dislikes}(x,y) \\ & \mathsf{dislikes}(x,z) \land \mathsf{Dish}(y) \land \mathsf{contains}(y,z) \to \mathsf{dislikes}(x,y) \\ & \mathsf{orderedDish}(x,y) \land \mathsf{dislikes}(x,y) \to \mathsf{Unhappy}(x) \end{split}$$





ThaiCurry $\sqsubseteq \exists$ contains.{peanutOil} $\top \sqsubseteq \forall$ orderedDish.Dish

 $\begin{array}{l} \text{NutAllergic(x)} \land \text{NutProduct(y)} \rightarrow \text{dislikes(x,y)} \\ \text{dislikes(x,z)} \land \text{Dish(y)} \land \text{contains(y,z)} \rightarrow \text{dislikes(x,y)} \\ \text{orderedDish(x,y)} \land \text{dislikes(x,y)} \rightarrow \text{Unhappy(x)} \end{array}$

Conclusions: dislikes(sebastian,peanutOil)





ThaiCurry ⊑ ∃contains.{peanutOil}

 $\top \sqsubseteq \forall orderedDish.Dish$

orderedDish rdfs:range Dish.

$$\begin{split} & \mathsf{NutAllergic}(x) \land \mathsf{NutProduct}(y) \to \mathsf{dislikes}(x,y) \\ & \mathsf{dislikes}(x,z) \land \mathsf{Dish}(y) \land \mathsf{contains}(y,z) \to \mathsf{dislikes}(x,y) \\ & \mathsf{orderedDish}(x,y) \land \mathsf{dislikes}(x,y) \to \mathsf{Unhappy}(x) \end{split}$$

Conclusions: dislikes(sebastian,peanutOil) orderedDish(sebastian,y_s) ThaiCurry(y_s) Dish(y_s)

inter 2010 – Pascal Hitzler



ThaiCurry ⊑ ∃contains.{peanutOil} ⊤ ⊑ ∀orderedDish.Dish

$$\begin{split} & \mathsf{NutAllergic}(x) \land \mathsf{NutProduct}(y) \to \mathsf{dislikes}(x,y) \\ & \mathsf{dislikes}(x,z) \land \mathsf{Dish}(y) \land \mathsf{contains}(y,z) \to \mathsf{dislikes}(x,y) \\ & \mathsf{orderedDish}(x,y) \land \mathsf{dislikes}(x,y) \to \mathsf{Unhappy}(x) \end{split}$$

Conclusions: dislikes(sebastian,peanutOil) orderedDish(sebastian,y_s) ThaiCurry(y_s)

Dish(y_s)

contains(y_s,peanutOil)

SWRL example



Does not work under DL-safety!

NutAllergic(sebastian) NutProduct(peanutOil) ∃orderedDish.ThaiCurry(sebastian)

ThaiCurry ⊑ ∃contains.{peanutOil} ⊤ ⊑ ∀orderedDish.Dish

$$\begin{split} & \mathsf{NutAllergic}(x) \land \mathsf{NutProduct}(y) \to \mathsf{dislikes}(x,y) \\ & \mathsf{dislikes}(x,z) \land \mathsf{Dish}(y) \land \mathsf{contains}(y,z) \to \mathsf{dislikes}(x,y) \\ & \mathsf{orderedDish}(x,y) \land \mathsf{dislikes}(x,y) \to \mathsf{Unhappy}(x) \end{split}$$

Conclusions: **dislikes(sebastian,peanutOil)** orderedDish(sebastian,y_s) ThaiCurry(y_s) **Dish(y_s)**

contains(y_s,peanutOil) dislikes(sebastian,y_s)



ThaiCurry ⊑ ∃contains.{peanutOil} ⊤ ⊑ ∀orderedDish.Dish

$$\begin{split} & \mathsf{NutAllergic}(x) \land \mathsf{NutProduct}(y) \to \mathsf{dislikes}(x,y) \\ & \mathsf{dislikes}(x,z) \land \mathsf{Dish}(y) \land \mathsf{contains}(y,z) \to \mathsf{dislikes}(x,y) \\ & \mathsf{orderedDish}(x,y) \land \mathsf{dislikes}(x,y) \to \mathsf{Unhappy}(x) \end{split}$$

Conclusions: dislikes(sebastian,peanutOil) orderedDish(sebastian,y_s) ThaiCurry(y_s) Dish(y_s)

contains(y_s,peanutOil) dislikes(sebastian,y_s) Unhappy(sebastian)

inter 2010 – Pascal Hitzler



ThaiCurry ⊑ ∃contains.{peanutOil} ⊤ ⊑ ∀orderedDish.Dish

$$\begin{split} & \mathsf{NutAllergic}(x) \land \mathsf{NutProduct}(y) \to \mathsf{dislikes}(x,y) \\ & \mathsf{dislikes}(x,z) \land \mathsf{Dish}(y) \land \mathsf{contains}(y,z) \to \mathsf{dislikes}(x,y) \\ & \mathsf{orderedDish}(x,y) \land \mathsf{dislikes}(x,y) \to \mathsf{Unhappy}(x) \end{split}$$

Conclusion: Unhappy(sebastian)





- SWRL and DL-safe SWRL are essentially based on the same style of model-theoretic semantics.
- If we want to deal with inconsistencies, uncertainty, or default reasoning, we have to modify the semantic approach.

- How to modify a semantics?
 - Redefine the notion of *model*!



Sophisticated Application Needs







Paraconsistent Reasoning

€ кпо.€.sis

• Modification: Use four truth values instead of two.

{true, false} \rightarrow {true, false, none, both}

- Idea: "both" captures inconsistency.
- Unicorn(beauty)
 Unicorn ⊑ Fictitious
 Unicorn ⊑ Animal
 Animal ⊑ ¬Fictitious

would, e.g., result in the truth value "both" for Fictitious(beauty).

• Problems: Paraconsistency or bugfixing? Which of various related approaches to take? How well does it work in practice?



Uncertainty Reasoning



- Modification: Use the real unit interval as set of truth values.
- 0 is interpreted as "false"
- 1 is interpreted as "true"
- Define how to combine them. E.g.,

HighQuality(a)is 0.7 trueExpensive(a)is 0.8 trueHighQuality □ Expensive □ Buyable

how "much" true is Buyable(a)?

• Problems: Different choices for combination. Does it match the intuition? Is this probabilistic or fuzzy? How reliable are the values? And it's computationally (much) more expensive.





- Thinkpads "normally" run Windows.
- I.e., this is the default assumption (to be assumed unless there is evidence to the contrary).

 Thinkpad ⊑ ∀runsOS.WindowsOS Thinkpad(myThinkpad) Thinkpad(yourThinkpad) runsOS(yourThinkpad,linux) ¬WindowsOS(linux)

is contradictory. How do we capture the default?





 Thinkpad ⊑ ∀runsOS.WindowsOS ⊔ ExceptionThing Thinkpad(myThinkpad) Thinkpad(yourThinkpad) runsOS(yourThinkpad,linux) ¬WindowsOS(linux)

+ a semantics which "minimizes" ExceptionThing. IAW, something is only in ExceptionThing if it is necessarily contained in it (e.g., to avoid a contradiction).

- This idea is called *circumscription* and is due to John McCarthy [1980] (not for DLs, obviously).
 There exist other approaches which accomplish the same thing in other ways.
- Problem: Computationally very expensive.





 Thinkpad ⊑ ∀runsOS.WindowsOS ⊔ ExceptionThing Thinkpad(myThinkpad) Thinkpad(yourThinkpad) runsOS(yourThinkpad,linux) ¬WindowsOS(linux)

+ a semantics which "minimizes" ExceptionThing. IAW, something is only in ExceptionThing if it is necessarily contained in it (e.g., to avoid a contradiction).

From all models I of the KB, select those models, for which ExceptionThing^I is minimal.
 Take these as the *circumscribed* models.
 Define logical consequence as usual.



Today: Wrap-Up



- Main Messages
- What Is Semantics Revisited
- OWL At Its Expressive Limits
- A Very Personal Semantic Web History
- Making OWL Fit For Practice



A Very Personal Semantic Web History

- 2002: Growing Semantic Web Hype (I wasn't there)
- 2004: Will it every work?
- 2006: It's probably not going to work.
- 2008: Industry is catching on and RDF will work. But OWL won't.
- 2010: Many major IT companies' R&D departments investigate OWL or even have their own OWL reasoner.



Today: Wrap-Up



- Main Messages
- What Is Semantics Revisited
- OWL At Its Expressive Limits
- A Very Personal Semantic Web History
- Making OWL Fit For Practice



Making OWL Fit For Practice



- The use of formal semantics for RDF and OWL still hasn't produced prominent applications with clear-cut added value compared with other methods/technologies.
 [More precisely: Such a thing hasn't been made public.]
- Hindrances:
 - Scalability of reasoning isn't very good (yet).
 - Few people can really model well in OWL.
 - High-quality ontologies are expensive to produce.
 - Real-life data often isn't clean enough for reasoning.



Making OWL Fit For Practice



- Researchers have to work on:
 - Scalability, including alternative reasoning methods.
 Don't get fixed on soundness/completeness/decidability.
 - Dissemination and education.
 - Methods for making real-life data fit for formal semantics.
 - Developing clear-cut use cases for formal semantics.

It is essential, to leave the ivory tower!





- Pascal Hitzler, Towards Reasoning Pragmatics. In: Krzysztof Janowicz, Martin Raubal, Sergei Levashkin (Eds.), GeoSpatial Semantics, Third International Conference, GeoS 2009, Mexico City, Mexico, December 3-4, 2009. Proceedings. Lecture Notes in Computer Science Vol. 5892 Springer 2009, pp. 9-25.
- Prateek Jain, Pascal Hitzler, Peter Z. Yeh, Kunal Verma, Amit P. Sheth, Linked Data is Merely More Data. Proceedings of Linked AI: AAAI Spring Symposium "Linked Data Meets Artificial Intelligence", March 22-24 2010. To appear.
- Markus Krötzsch, Sebastian Rudolph, Pascal Hitzler, ELP: Tractable Rules for OWL 2. In: Amit Sheth, Steffen Staab, Mike Dean, Massimo Paolucci, Diana Maynard, Timothy Finin, Krishnaprasad Thirunarayan (eds.), The Semantic Web - ISWC 2008, 7th International Semantic Web Conference. Springer Lecture Notes in Computer Science Vol. 5318, 2008, pp. 649-664.
- Markus Krötzsch, Sebastian Rudolph, Pascal Hitzler, Description Logic Rules. In: Malik Ghallab, Constantine D. Spyropoulos, Nikos Fakotakis, Nikos Avouris (eds.), Proceedings of the 18th European Conference on Artificial Intelligence, ECAI2008, Patras, Greece, July 2008. IOS Press, 2008, pp. 80-84





- Stephan Grimm, Pascal Hitzler, A Preferential Tableaux Calculus for Circumscriptive ALCO. In: Axel Polleres, Terrance Swift (Eds.), Web Reasoning and Rule Systems, Third International Conference, RR 2009, Chantilly, VA, USA, October 20009, Proceedings. Lecture Notes in Computer Science Vol. 5837, Springer, pp. 40-54.
- Yue Ma, Pascal Hitzler, Paraconsistent reasoning for OWL 2. In: Axel Polleres, Terrance Swift (Eds.), Web Reasoning and Rule Systems, Third International Conference, RR 2009, Chantilly, VA, USA, October 20009, Proceedings. Lecture Notes in Computer Science Vol. 5837, Springer, pp. 197-211.
- Sebastian Rudolph, Tuvshintur Tserendorj, Pascal Hitzler, What Is Approximate Reasoning? In: D. Calvanese, G. Lausen (eds.), Web Reasoning and Rule Systems, Second International Conference, RR 2008, Karlsruhe, Germany, October/November 2008. Springer Lecture Notes in Computer Science Vol. 5341, 2008, pp. 150-164.

